# Optimal allocation of buffer times to increase train schedule robustness 

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1.0


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- Introduction and motivation
- Knapsack problem approach
- Parameter computation
- Case study and results


## Capacity4Rail WP3.2

- Simulation and models to evaluate enhanced capacity
- The aim of this task is to evaluate existing tools for their suitability to assess and improve capacity utilization
- "Capacity depends on the way it is utilised" (UIC 406)
- Timetabling (and traffic control) determine the way capacity is utilised

Timetabling \& Traffic control


## Timetabling - C4R Perspective



## Research question

Robust timetabling enable more trains to run


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## Delay propagation - microscopic



## Delay propagation - macroscopic



## Solution - Buffer times




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## Another problem



## Existing solutions

## Robust optimisation:

+ Attacks general problems
- Very difficult to solve

Domain knowledge used to relax the problem:

+ Emma: Critical points
+ Fahimeh: Travel time dependent buffering


## Produce a general solution using the domain knowledge

## Problem definition

- Input - Timetable A
- Number of trains
- Scheduled running and dwell times
- Fixed train sequence
- Time window constraint
- Output - Timetable B
- All properties of Timeble A are kept
- Buffer times (re)distributed to increase robustness


## Knapsack problem (1/2)

- Hikers wants to go on a trip

- The backpack is small, no more than 10 kg of things in the bag
- He has prepared a list of items that he would like to bring on a trip
- Water, bread, cans, maps \& compass, laptop, trousers, jacket, socks \& underwear, knife and cutlery, sweater, tent, sleeping bag


## Knapsack problem (2/2)

| Item | Weight <br> [kg] | Utility: <br> 1 (not useful) to 10 (very useful) |
| :--- | :--- | :--- |
| Cans | 2.2 | 7 |
| Water | 2 | 4 |
| Tent | 3.5 | 8 |
| Food | 3 | 8 |
| Jacket | 0.5 | 7 |
| Maps \& compass | 0.1 | 10 |
| Sleeping bag | 0.8 | 9 |
| Laptop | 1.5 | 3 |
| Trousers | 0.3 | 6 |
| Socks \& underwear | 0.2 | 9 |
| Knife \& cutlery | 0.5 | 9 |
| Sweater | 0.5 | 5 |

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## Knapsack problem for buffer times (1/2)

Timetable compression UIC 406 -ish


## Knapsack problem for buffer times (1/2)



## Knapsack problem for buffer times (2/2)



- How to coordinate multiple sections?
- How to prioritize items (candidates)?
- Marginal profit: is the second minute (time unit) of buffer as valuable as the first? How about the third?


## Multidimensional Knapsack Problem



Figure 3: Illustrative example for the knapsack capacity


Figure 4: Illustrative example for the knapsack capacity

## Prioritisation

- Efficient graph algorithms can be used to compute for each candidate:
- 1. Delay impact (I): if the candidate is delayed for $D$, how many events will have secondary delay?
- 2. Delay sensitivity (S): how many other events can be delayed for $D$ so that it propagates to the candidate?
- The bigger I and S, the bigger the profit for including the candidate!

) second train without scheduled stop

e) overtaking


d) first train without scheduled stop



## Marginal profit

- Marginal profit from including an additional minute depends on the number of already included minutes of the same buffer



## Case study



## Case study



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## Case study



## Experimental setup

- 3 schedules generated by using different parameter setup
- 500 hundred primary delay scenarios generated
- All departured events are delayed with a uniform distribution upto 10 minutes
- On average 28 events have primary delay
- Total primary delay 150.14 min on average
- Deterministic delay propagation algorithm computed secondary delays in each scenario for each timetable ( $500 \times 4$ experiments in total)
- Upto $11 \%$ decrease in seconary delay


## Results

|  | Total delay <br> $[\mathrm{min}]$ | Average <br> delay per <br> event [min] | Delay per <br> 1 min prim. <br> $[\mathrm{min}]$ | Delay per <br> init. delayed <br> event [min] |
| :--- | :---: | :---: | :---: | :---: |
| Original | 1146.70 | 8.49 | 8.87 | 40.95 |
| TB 0-1 | 1034.20 | 7.66 | 7.99 | 36.94 |
| TB Bounded | 1033.80 | 7.65 | 7.98 | 36.92 |
| TTB | 1017.20 | $\mathbf{7 . 5 3}$ | $\mathbf{7 . 8 4}$ | $\mathbf{3 6 . 3 2}$ |



- In upto $87 \%$ cases, original timetable performes worse


## Next steps

- Prioritisation of buffering base don historical data
- Compuational experinements on networks
- More details about the approach available soon:
- Jovanovic P., Kecman P., Bojovic N., Mandic D. Optimal allocation of buffer times to increase schedule robustness. European Journal of Oprerations Research (to appear soon)


## Thank you for your attention

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